

Remnants from the Bookshelf

cropped up again in jww Silvia Steila

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White spots within a basically known area

Γ_0

Bachmann-Howard ordinal

$\Pi_1^1\text{-CA} + (\text{BI})$

$\Delta_2^1\text{-CA} + (\text{BI})$

$\Pi_2^1\text{-CA} + (\text{BI})$

The general framework

The languages \mathcal{L}_1 and \mathcal{L}^*

- \mathcal{L}_1 : a standard language of first order arithmetic with a constant \bar{m} for every natural number m and an n -ary relation symbol $R_{\mathcal{Z}}$ for every n -ary primitive recursive relation \mathcal{Z} ; we write 0 for $\bar{0}$.
- $\mathcal{L}^* := \mathcal{L}_1(N, S, Ad, \in)$
 - ▶ N a constant for the set of the natural numbers,
 - ▶ S a unary relation symbol to say that an object is a set.
 - ▶ Ad a unary relation symbol to say that an object is an admissible set.
- Δ_0 , Σ_n , Π_n , Σ , and Π formulas of \mathcal{L}^* defined as usual.
- A^b is the relativization of formula A to set b .

Systems of basic set theory BS^0 and BS

Number-theoretic axioms of BS^0

(N.1) A^N for every closed axiom A of PA.

(N.2) $0 \in a \wedge (\forall x, y \in N)(x \in a \wedge R_S(x, y) \rightarrow y \in a) \rightarrow N \subseteq a$.

Ontological axioms of BS^0 .

(O.1) $S(a) \leftrightarrow a \notin N$.

(O.2) $a \in N \rightarrow b \notin a$.

(O.3) $0 \in N$.

(O.4) $R_Z(a_1, \dots, a_n) \rightarrow a_1, \dots, a_n \in N$.

Set-theoretic axioms of BS^0 .

(S.1) Pair: $\exists x(a \in x \wedge b \in x)$.

(S.2) Union: $\exists x(\forall y \in a)(\forall z \in y)(z \in x)$.

(S.3) Δ_0 separation (Δ_0 -Sep): for all Δ_0 formulas $\varphi[y]$,

$$\exists x(x = \{y \in a : \varphi[y]\}).$$

$BS := BS^0$ plus full induction on N and full \in -induction, i.e.

$$\forall x((\forall y \in x)\varphi[y] \rightarrow \varphi[x]) \rightarrow \forall x\varphi[x]$$

for all formulas $\varphi[x]$ of \mathcal{L}^* .

Theories for admissible sets

The schema of Δ_0 collection

For all Δ_0 formulas $\varphi[x, y]$ of \mathcal{L}^* :

$$(\forall x \in a)\exists y\varphi[x, y] \rightarrow \exists z(\forall x \in a)(\exists y \in z)\varphi[x, y]. \quad (\Delta_0\text{-Col})$$

KPU⁰ and KPU

$$\text{KPU}^0 := \text{BS}^0 + (\Delta_0\text{-Col}),$$

$$\text{KPU} := \text{BS} + (\Delta_0\text{-Col}).$$

Theorem (Jä)

$$\text{KPU}^0 \equiv \text{PA} \quad \text{and} \quad \text{KPU} \equiv \text{ID}_1.$$

Adding admissible sets

Ad-axioms

(*Ad.1*) $Ad(d) \rightarrow N \in d \wedge Tran[d]$.

(*Ad.2*) $Ad(d_1) \wedge Ad(d_2) \rightarrow d_1 \in d_2 \vee d_1 = d_2 \vee d_2 \in d_1$.

(*Ad.3*) For any closed instance of an axiom φ of KPu,

$$Ad(d) \rightarrow \varphi^d.$$

Remark. The *Ad*-axioms do not imply the existence of admissible sets. However, this is achieved by the limit axiom (Lim):

$$\forall x \exists y (x \in y \wedge Ad(y)).$$

KPI⁰, KPI, KPi⁰, KPi

$$\text{KPI}^0 := \text{BS}^0 + \text{Ad-axioms} + (\text{Lim}),$$

$$\text{KPI} := \text{BS} + \text{Ad-axioms} + (\text{Lim}),$$

$$\text{KPi}^0 := \text{KPU}^0 + \text{Ad-axioms} + (\text{Lim}),$$

$$\text{KPi} := \text{KPU} + \text{Ad-axioms} + (\text{Lim}),$$

Theorem (Jä)

$$\text{KPI}^0 \equiv \text{KPi}^0 \equiv \text{ATR}_0,$$

$$\text{KPI} \equiv \Pi_1^1\text{-CA} + (\text{BI}),$$

$$\text{KPi} \equiv \Delta_2^1\text{-CA} + (\text{BI}).$$

The role of Δ_1 separation

The schema of Δ_1 separation

For any Σ_1 formula $\varphi[x]$ and Π_1 formula $\psi[x]$ of \mathcal{L}^* ,

$$(\forall x \in a)(\varphi[x] \leftrightarrow \psi[x]) \rightarrow \exists y(y = \{x \in a : \varphi[x]\}). \quad (\Delta_1\text{-Sep})$$

We know:

- (i) $(\Delta_1\text{-Sep})$ is provable in KPU^0 .
- (ii) Jensen, cf. Barwise: α is admissible if and only if α is a limit ordinal and L_α satisfies $(\Delta_1\text{-Sep})$.

Conjecture/Theorem

$$\text{BS}^0 + (\Delta_1\text{-Sep}) \equiv \Delta_1^1\text{-CA}_0 \quad \text{and} \quad \text{BS} + (\Delta_1\text{-Sep}) \equiv \Delta_1^1\text{-CA}.$$

Question

- Is there a natural way to formalize something like $(V=L)$ in the theory $BS + (\Delta_1\text{-Sep})$?
- And if so, does it affect the proof-theoretic strength of this theory?

Σ and Π reduction

Definition

Let \mathfrak{F} be a collection of formulas of \mathcal{L}^* and $\neg\mathfrak{F}$ the collection of its duals. Then the axioms schema (\mathfrak{F} -Red) of \mathfrak{F} reduction consists of all formulas

$$(\forall x \in a)(\varphi[x] \rightarrow \psi[x]) \rightarrow \exists y(\{x \in a : \varphi[x]\} \subseteq y \subseteq \{x \in a : \psi[x]\}),$$

where $\varphi[x]$ is from $\neg\mathfrak{F}$ and $\psi[x]$ from \mathfrak{F} .

Conjecture/Theorem

- ① $BS^0 + (\Sigma\text{-Red}) \equiv \Sigma_1^1\text{-AC}_0$ and $BS + (\Sigma\text{-Red}) \equiv \Sigma_1^1\text{-AC}$.
- ② $BS^0 + (\Pi\text{-Red}) \equiv \text{ATR}_0$ and $BS + (\Pi\text{-Red}) \equiv \text{ATR}$.

Question

Clearly, KPu^0 proves $(\Sigma\text{-Red})$. But what can we say about

$$KPu^0 + (\Pi\text{-Red}) \quad \text{and} \quad KPu + (\Pi\text{-Red})?$$

It is clear that

$$KPu^0 + (\Pi\text{-Red}) \subseteq KPu^0 + (\Sigma_1\text{-Sep}),$$

$$KPu + (\Pi\text{-Red}) \subseteq KPu + (\Sigma_1\text{-Sep}).$$

Adding principles of second order arithmetic

Canonical embedding of \mathcal{L}_2 into \mathcal{L}^* .

Language \mathcal{L}_2 of second order arithmetic embedded into \mathcal{L}^* by translating

$$\begin{aligned} \exists n(\dots) &\text{ into } (\exists n \in \mathbb{N})(\dots) \quad \text{and} \quad \forall n(\dots) &\text{ into } (\forall n \in \mathbb{N})(\dots), \\ \exists X(\dots) &\text{ into } (\exists x \subseteq \mathbb{N})(\dots) \quad \text{and} \quad \forall X(\dots) &\text{ into } (\forall x \subseteq \mathbb{N})(\dots). \end{aligned}$$

Theorem

- 1 $\text{KP}_u \equiv \text{ID}_1 \equiv \text{ACA} + (\Pi_1^1\text{-CA})^-$.
- 2 $\text{KP}_i \equiv \Delta_2^1\text{-CA} + (\text{BI})$.

Question

What is the strength of $\text{KP}_u + (\Pi_1^1\text{-CA})$?

Arithmetic operator forms

Arithmetic formulas (in their translation into \mathcal{L}^*) of the form

$$\mathfrak{A}[X^+, n],$$

possibly with additional set and number parameters. We set

$$Fix_{\mathfrak{A}}[\mathbb{N}, x] := x \subseteq \mathbb{N} \wedge (\forall n \in \mathbb{N})(n \in x \leftrightarrow \mathfrak{A}[x, n]).$$

Arithmetic fixed point axioms

Let $\mathfrak{A}[X^+, n]$ be an arithmetic operator form.

$$\exists x Fix_{\mathfrak{A}}[\mathbb{N}, x], \quad (\Pi_{\infty}^0\text{-FP})$$

$$\exists x (Fix_{\mathfrak{A}}[\mathbb{N}, x] \wedge \forall y (Fix_{\mathfrak{A}}[\mathbb{N}, y] \rightarrow x \subseteq y)), \quad (\Pi_{\infty}^0\text{-LFP})$$

Questions

- 1 Is $KPu^0 + (\Pi_\infty^0\text{-FP})$ proof-theoretically equivalent to ATR_0 ?
- 2 What is the proof-theoretic strength of $KPu + (\Pi_\infty^0\text{-FP})$?
- 3 What is the proof-theoretic strength of $KPu + (\Pi_\infty^0\text{-LFP})$?

Subset-bounded separations

$\Pi_1^{\mathcal{P}}$ separation ($\Pi_1^{\mathcal{P}}$ -Sep)

For any Δ_0 formula $\varphi[x, y]$ and any a :

$$\exists z (z = \{x \in a : (\forall y \subseteq a) \varphi[x, y]\}).$$

$\Pi_1^{\mathcal{P}}(\Delta_1)$ separation ($\Pi_1^{\mathcal{P}}(\Delta_1)$ -Sep)

For every Σ_1 formula $\varphi[x, y]$, every Π formula $\psi[x, y]$, and any a :

$$\forall x, y (\varphi[x, y] \leftrightarrow \psi[x, y]) \rightarrow \exists z (z = \{x \in a : (\forall y \subseteq a) \varphi[x, y]\}).$$

Questions

- 1 What is the exact relationship between $KP_u + (\Pi_1^1\text{-CA})$ and $KP_u + (\Pi_1^{\mathcal{P}}\text{-Sep})$?
- 2 What is the exact relationship between $KP_u + (\Pi_1^{\mathcal{P}}\text{-Sep})$ and $KP_u + (\Delta_0\text{-LFP})$?
- 3 Is $(\Delta_2^1\text{-CA})$ provable in $KP_u + (\Pi_1^{\mathcal{P}}(\Delta_1)\text{-Sep})$?

Thank you for your attention!