# How large are proper classes?

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# ABM

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# NGB

The theory NGB is formulated in a two-sorted language and consists of the following axioms:

- extensionality, pair, union, powerset, infinity for sets,
- Extensionality, Foundation for classes,
- Class Comprehension Schema: i.e, for every formula φ containing no quantifiers over classes there exists a class C such that

$$\forall x(\varphi[x] \leftrightarrow x \in C)$$

▶ Limitation of Size: i.e, for every proper class *C* there is a bijection between *C* and the class *V* of all sets.

• Let  $\mathcal{L}^c$  be the extension of  $\mathcal{L}$  with countably many class variables.

► The atomic formulas comprise the ones of *L* and all expression of the form "a ∈ C".

• An  $\mathcal{L}^c$  formula is elementary if it contains no class quantifiers.

 Δ<sup>c</sup><sub>n</sub>, Σ<sup>c</sup><sub>n</sub> and Π<sup>c</sup><sub>n</sub> are defined as usual, but permitting subformulas of the form "a ∈ C".

## **KP**<sup>c</sup>

The theory  $KP^c$  is formulated in  $\mathcal{L}^c$  and consists of the following axioms:

- extensionality, pair, union, infinity,
- Δ<sup>c</sup><sub>0</sub>-Separation: i.e, for every Δ<sup>c</sup><sub>0</sub> formula φ in which x is not free and any set a,

$$\exists x (x = \{y \in a : \varphi[y]\})$$

•  $\Delta_0^c$ -Collection: i.e, for every  $\Delta_0^c$  formula  $\varphi$  and any set a,

$$\forall x \in a \exists y \varphi[x, y] \to \exists b \forall x \in a \exists y \in b \varphi[x, y]$$

► Δ<sup>c</sup><sub>1</sub>-Comprehension</sub>: i.e, for every Σ<sup>c</sup><sub>1</sub> formula φ and every Π<sup>c</sup><sub>1</sub> formula ψ,

$$\forall x(\varphi[x] \leftrightarrow \psi[x]) \rightarrow \exists X \forall x(x \in X \leftrightarrow \varphi[x])$$

• Elementary  $\in$ -induction: i.e, for every elementary formula  $\varphi$ ,

$$\forall x ((\forall y \in x \varphi[y]) \rightarrow \varphi[x]) \rightarrow \forall x \varphi[x]$$

# Motivations: ... last ABM



### **Operators**

We call a class an operator if all its elements are ordered pairs and it is right-unique (i.e. functional).

▶ We use *F* to denote operators.

• Given an operator F and a set a we write Mon[F, a] for:

$$\forall x(F(x) \subseteq a) \land \forall x, y(x \subseteq y \rightarrow F(x) \subseteq F(y)).$$

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## Least fixed point statements



### LFP

$$\mathsf{Mon}[F,a] \to \exists x (F(x) = x \land \forall y (F(y) = y \to x \subseteq y)$$

# Separation

### $\Sigma_1^c$ -separation

For every  $\Sigma_1^c$  formula  $\varphi$  in which x is not free and any set a,

$$\exists x(x = \{y \in a : \varphi[y]\}).$$

# $\mathsf{SBS}\;(\sim \mathsf{\Pi}^{\mathcal{P}}_1(\Delta^c_1)\text{-}\mathsf{Sep})$

For every  $\Delta_1^c$  formula  $\varphi$  and sets *a* and *b*,

 $\exists z(z = \{x \in a : \exists y \subseteq b(\varphi[x, y])\})$ 

Fixed point principles in  $KP^{c} + (V=L)$ 



If we add to our theory the Axiom of Limitation of Size:

• we have a global well-ordering of V,

all our principles are equivalent,

But... I am not able to prove the consistency of: KP<sup>c</sup> + FP + Limitation of size, from the consistency of KP<sup>c</sup> + FP.

What does it happen if we consider something weaker than a bijection?

# Injections from ordinals to reals

### Proposition

Assume that there are no injections from Ord to  $\mathcal{P}(\omega).$  Then MI hold!

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# Injections from ordinals to reals

### Proposition

Assume that there are no injections from Ord to  $\mathcal{P}(\omega)$ . Then MI hold!

#### Question

And if there is an injection from Ord to  $\mathcal{P}(\omega)$ ?

## Injections from reals to ordinals

### Proposition

Assume that there is an injection from  $\mathcal{P}(\omega)$  to Ord. Then BPI implies MI.

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## Injections from reals to ordinals

#### Proposition

Assume that there is an injection from  $\mathcal{P}(\omega)$  to Ord. Then BPI implies MI.

#### Question

Assume that there are no injections from  $\mathcal{P}(\omega)$  to Ord... BPI holds.

## Surjections from ordinals to reals

#### Proposition

Assume that there is a surjection from Ord to  $\mathcal{P}(\omega)$ . Then there exists a strong well ordering of  $\mathcal{P}(\omega)$ .

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#### Proposition

Assume that there is a surjection from Ord to  $\mathcal{P}(\omega)$ . Then there exists a strong well ordering of  $\mathcal{P}(\omega)$ .

#### Question

Which is the strength of the statement: "For every class C, there exists either an injection from C to the ordinals or a surjection from the ordinals to C"?

#### Theorem

Assume that there exists a cofinal map  $F : \mathcal{P}(\omega) \to \text{Ord.}$  Then SBS implies  $\Sigma_1^c$ -Separation for ordinals.

- Given  $\varphi$  we want to show that  $\{x \in \omega : \exists \alpha \varphi[\alpha, x]\}$  is a set.
- ► By using *F*:

$$\exists \alpha \varphi[x, \alpha] \iff \exists y \subseteq \omega(\exists \alpha < F(y)(\varphi[x, \alpha])).$$

- The formula " $\exists \alpha < F(y)(\varphi[x, \alpha])$ " is  $\Delta^c$ .
- By applying SBS we get the thesis.

Let CM be the statement: there exists a cofinal map  $F : \mathcal{P}(\omega) \to \text{Ord}$ .

• 
$$L \models (\mathsf{CM} \lor (\mathcal{P}(\omega) \text{ is a set})).$$

Axiom Beta does not imply CM.

CM does not imply Axiom Beta.

CM does not imply that every the least fixed point of any arithmetical operator is Δ<sup>c</sup>-definable.

What about the negation of CM?



#### Theorem

Assume that there are no cofinal maps from the reals to the ordinals. Then  $\Pi_1$ -Reduction for ordinals holds.

#### $\Pi_1$ -Reduction for ordinals

Let  $\varphi$  and  $\psi$  be two  $\Delta_0$  formulas such that

$$\forall x \in \omega (\exists \alpha \varphi[x, \alpha] \implies \forall \alpha \psi[x, \alpha]).$$

there exists a set z such that

 $\{x \in \omega : \exists \alpha \varphi[x, \alpha]\} \subseteq z \subseteq \{x \in \omega : \forall \alpha \psi[x, \alpha]\}.$ 

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• Assume that we have a set  $\omega$  and two  $\Delta$  formulas  $\varphi$  and  $\psi$  such that

$$\forall x \in \omega (\exists \alpha \varphi[x, \alpha] \implies \forall \alpha \psi[x, \alpha])$$

and  $\Pi_1$ -Reduction for them does not hold.

We derive

$$\forall z \subseteq \omega \exists x \in \omega \exists \alpha ((\varphi[x, \alpha] \land x \notin z) \lor (x \in z \land \neg \psi[x, \alpha]))$$

• Define the following operator  $F : \mathcal{P}(\omega) \to \text{Ord}$ .

$$F(z) = \mu \alpha (\exists x (\varphi[x, \alpha] \land x \notin z) \lor (x \in z \land \neg \psi[x, \alpha])).$$

• There exists  $\beta$  such that

$$\forall z \subseteq \omega \exists x \in \omega \exists \alpha \in \beta ((\varphi[x, \alpha] \land x \notin z) \lor (x \in z \land \neg \psi[x, \alpha]))$$

Define the set

$$\{x \in \omega : \exists \alpha < \beta \varphi[x, \alpha]\}.$$

and derive a contradiction.

Moreover:

- SBS implies  $\Pi_1$ -Reduction for ordinals.
- The Axiom of Powerset implies  $\neg$ CM.

► ¬CM does not imply Axiom Beta.

### Question

Which is the strength of Π<sub>1</sub>-Reduction for ordinals?

Does Axiom Beta imply ¬CM?

Moreover:

- SBS implies  $\Pi_1$ -Reduction for ordinals.
- The Axiom of Powerset implies  $\neg$ CM.

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### Question

- Which is the strength of Π<sub>1</sub>-Reduction for ordinals?
- Does Axiom Beta imply ¬CM?

Thank you!

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